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Incentive Mechanisms for Spatio-temporal Tasks in Mobile Crowdsensing

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Mobile Crowd Sensing

light sensor

Digital Compass

Accelerometer



Camera

GPS

proximity sensor Microphone Gyroscope

traffic monitoring noise monitoring pollution monitoring platform garbage classification

Incentive mechanisms for mobile crowd sensing





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compensate users' cost

help to achieve good service quality







Related Work

- Location dependent task
 - Y. Chon, et al., UbiComp, 2013.
 - Y. Feng, et al., INFOCOM, 2014.
- Time dependent task
 - J. Xu, et al., IEEE Trans. on Wireless Communications, 2015.
 - K. Han, et al., IEEE Trans. on Computers, 2016.
 - J. Xu, et al., Wireless Networks, 2017.
- Spatio-temporal tasks without overlap
 - Q. Li, et al., PerCom, 2013.
 - Z. Wang, et al., Computer Networks, 2018.





Spatio-temporal Tasks for Traffic Monitoring







- How to determine the value of sensing data provided by such users who can contribute for multiple tasks simultaneously?
- How to allocate the sensing time of the mobile users for their sensing areas?
- How to prevent the strategic behavior by submitting dishonest bidding price?





Location Sensitive Model

Location Sensitive Social Optimization (LSSO) problem



The LSSO problem is NP-hard since it is the Multi-set Multi-cover problem

Incentive <u>Mechanism for</u> <u>Location</u> <u>Sensitive</u> Model (MLS)

The effective average cost of user *i*

$$\frac{b_i}{\sum_{\tau_j \in \Gamma_i} \min\left\{t_i, t'^j\right\}}$$

Phase1: winner selection

while
$$\sum_{\tau_j \in \Gamma} t'^j \neq 0$$
 do
 $i \leftarrow \arg \min_{h \in U \setminus S} \frac{b_h}{\sum_{\tau_j \in \Gamma_h} \min\{t_h, t'^j\}};$
 $S \leftarrow S \cup \{i\};$
for each $\tau_j \in \Gamma_i$ do $t'^j \leftarrow t'^j - \min\{t_i, t'^j\};$
end

select the user with minimum effective average cost over the unselected user set as the winner until the winners' sensing time can meet the requirement of minimum sensing time of each task.

Incentive <u>Mechanism for</u> <u>Location</u> <u>Sensitive</u> Model (MLS)



Theoretical Analysis of MLS

Lemma 1. MLS is computationally efficient

 $O(n^3\varepsilon)$, ε is the maximum of overlaps of Aols

Lemma 2. MLS is individually rational.

Each winner will have a nonnegative utility while bidding its true cost.

Lemma 3. MLS is truthful

No user can improve its utility by submitting a false cost, no matter what others submit.

Lemma 4. MLS can approximate the optimal solution within a factor of H_K , where $K = \max_{i \in U} \sum_{\tau_i \in \Gamma_i} \min\{t_i, t^j\}, H_K = 1 + \frac{1}{2} + \dots + \frac{1}{K}$.



The *LISO* problem is NP-hard since the *LISO* problem is a generalization of the *LSSO* problem.

Incentive <u>Mechanism for</u> <u>Location</u> <u>Insensitive</u> Model (MLI)

The effective average cost of user *i*



Phase1: winner selection

while
$$\sum_{\tau_j \in \Gamma} t'^j \neq 0$$
 do
TSA \leftarrow Allocation $(\Gamma, U \setminus S, \mathbf{B}, \{t'^1, t'^2, \dots, t'^m\});$
 $i \leftarrow \arg \min_{h \in U \setminus S} \frac{b_h}{\sum_{\tau_j \in \Gamma_h} \sum_{k:sa_{h,k} \cap a^j \neq \emptyset} (tsa_{h,k} \cdot |(sa_{h,k}|))};$
 $S \leftarrow S \cup \{i\};$
for each $\tau_j \in \Gamma_i$ do $t'^j \leftarrow t'^j - \sum_{k:sa_{h,k} \cap a^j \neq \emptyset} tsa_{h,k};$
end

calculate the sensing time for each area by calling function Allocation(·) for each user

Incentive Mechanism for Location Insensitive Model (MLI)

calculate the sensing time allocation matrix

select the sensing areas with most overlaps of AoIs

select the task with minimum residual sensing time of all tasks overlapping the selected sensing area

calculate the minimum value of this two time

foreach $i \in U''$ do $t'_i \leftarrow t_i; SA'_i \leftarrow SA_i; \{\overline{t'^1}, \overline{t'^2}, \dots, \overline{t'^m}\} \leftarrow \mathcal{R};$ while $t'_i > 0$ and $SA'_i \neq \emptyset$ do $k \leftarrow \arg \max_{k:sa_{i,k'} \in SA'_i} |sa_{i,k'}|;$ $j \leftarrow \arg \min_{j':a^{j'} \cap sa_{i,k} \neq \emptyset} \overline{t'^{j'}};$ $tsa_{i,k} \leftarrow \min\{\overline{t'^j}, t'_i\};$ $t'_i \leftarrow t'_i - tsa_{ik};$ $SA'_i \leftarrow SA'_i \setminus \{sa_{i,k}\};$ for each $\tau_i \in \Gamma_i$ s.t. $a^j \cap sa_{i,k} \neq \emptyset$ do $\overline{t^{\prime j}} \leftarrow \overline{t^{\prime j}} - tsa_{i,k};$ end end end

Incentive <u>Mechanism for</u> <u>Location</u> <u>Insensitive</u> Model (MLI)



Theoretical Analysis of MLI

MLI is computationally efficient, individually rational, truthful, and H_K *approximate, where* $K = \max_{i \in U} \sum_{\tau_j \in \Gamma_i} \sum_{k:sa_{i,k} \cap a^j \neq \emptyset} (tsa_{i,k} \cdot |(sa_{i,k}|)).$

A Toy Example for Winner Selection of MLI



	Userl			User2	User3	User4	User5
ti	4			1	2	4	4
b_i	5			2	3	7	9
	tsa _{1,1}	tsa _{1,2}	tsa _{1,3}	tsa _{2,1}	tsa _{3,1}	tsa _{4,1}	tsa _{5,1}
Round 1	4	0	0	1	2	4	4
Winner							
Round 2				1	2	1	2
Winner							
Round 3				1		1	0
Winner							

Round 1:
$$t'^1 = 5, t'^2 = 6, S = \emptyset$$

 $\frac{b_1}{2tsa_{1,1}} = \frac{5}{8}, \frac{b_2}{tsa_{2,1}} = 1, \frac{b_3}{tsa_{3,1}} = \frac{3}{2}, \frac{b_4}{tsa_{4,1}} = \frac{7}{4}, \frac{b_5}{tsa_{5,1}} = \frac{9}{4}.$
user 1 wins

A Toy Example for Winner Selection of MLI



	Userl			User2	User3	User4	User5
ti	4			1	2	4	4
b_i	5			2	3	7	9
	tsa _{1,1}	tsa _{1,2}	tsa _{1,3}	tsa _{2,1}	tsa _{3,1}	tsa _{4,1}	tsa _{5,1}
Round 1	4	0	0	1	2	4	4
Winner							
Round 2				1	2	1	2
Winner							
Round 3				1		1	0
Winner							

Round 2:
$$t'^1 = 5 - 4 = 1$$
, $t'^2 = 6 - 4 = 2$, $S = \{1\}$
 $\frac{b_2}{tsa_{2,1}} = 2$, $\frac{b_3}{tsa_{3,1}} = \frac{3}{2}$, $\frac{b_4}{tsa_{4,1}} = 7$, $\frac{b_5}{tsa_{5,1}} = \frac{9}{2}$.
user 3 wins.

A Toy Example for Winner Selection of MLI



	Userl			User2	User3	User4	User5
ti	4			1	2	4	4
b_i	5			2	3	7	9
	tsa _{1,1}	tsa _{1,2}	tsa _{1,3}	tsa _{2,1}	tsa _{3,1}	tsa _{4,1}	tsa _{5,1}
Round 1	4	0	0	1	2	4	4
Winner							
Round 2				1	2	1	2
Winner							
Round 3				1		1	0
Winner							

Round 3:
$$t'^1 = 1, t'^2 = 2 - 2 = 0, S = \{1,3\}$$

 $\frac{b_2}{tsa_{2,1}} = 2, \frac{b_4}{tsa_{4,1}} = 7$. user 2 wins.
Thus $S = \{1,3,2\}$.

Performance Evaluation

Bench Mark Algorithms

- *MLS-GB* : greedily select the user with minimum bidding price as the winner in the location sensitive model
- *MLS-GC:* greedily select the user with maximal effective coverage as the winner in the location insensitive model
- *MLI-GB* : greedily select the user with minimum bidding price as the winner in the location sensitive model
- *MLI-GC:* greedily select the user with maximal effective coverage as the winner in the location insensitive model
- *ApproxMCS:* untruthful approximation algorithm [1] for maximizing the revenue of owner in mobile crowdsensing

[1]K. Han, C. Zhang, J. Luo, M. Hu, and B. Veeravalli, "Truthful Scheduling Mechanisms for Powering Mobile Crowdsensing," *IEEE Trans. on Computers*, vol.65, no.1, pp. 294-307, 2016.

Performance Evaluation

Datasets

- air pollution data [2] from the sites in Beijing
- T-Drive trajectory data [3] in Beijing
- contains trajectories of 10,357 taxis in Beijing
- regard the taxi trajectory between 14:30:29 and 15:00:29 as the active areas of the user

[2]http://beijingair.sinaapp.com/

[3]https://www.microsoft.com/en-us/research/publication/t-drive-trajectory-data-sample/



MLS outputs 22.3% and 5.3% less social cost than *MLS-GB* and *MLS-GC* on average, respectively. *MLI* outputs 33.6% and 7.8% less social cost than *MLI-GB* and *MLI-GC* on average, respectively.

Thank you! Q&A

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